Secure Sockets Layer

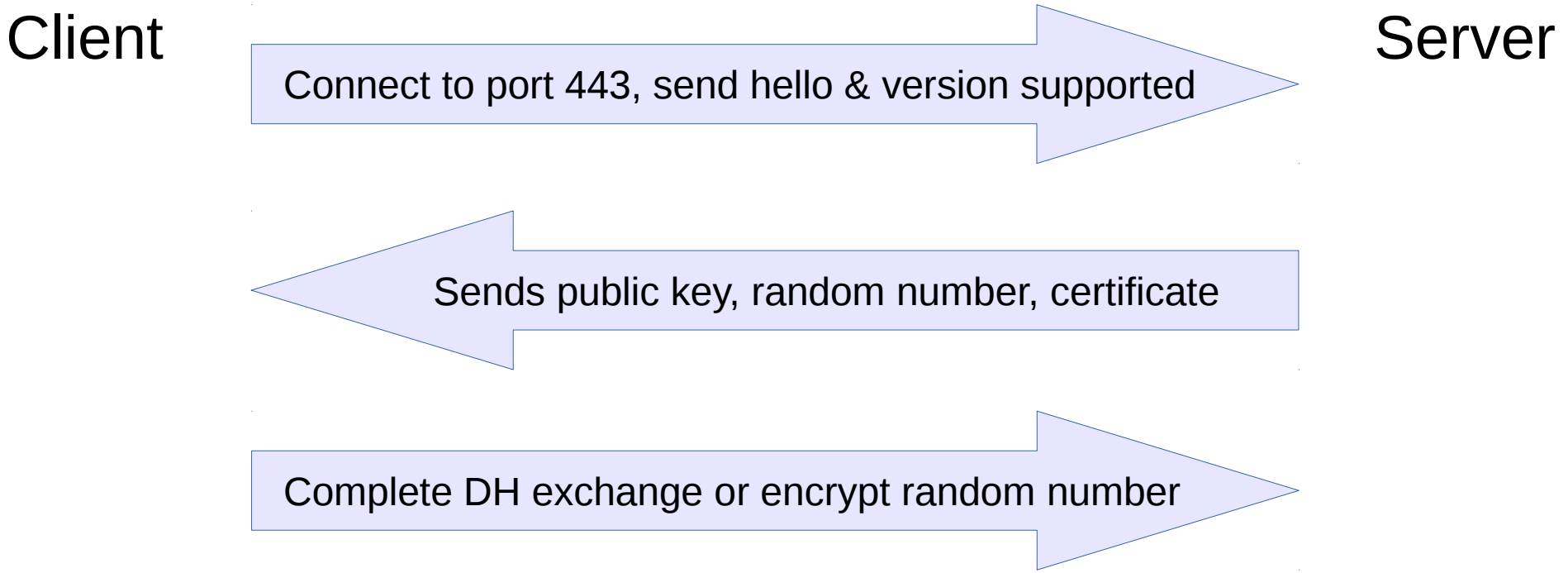
and

Transport Layer Security

Secure Sockets Layer

- Introduced by Netscape in 1995
- Initially a closed proprietary protocol
- SSL 3 protocol was published in 1996
- Standardized as Transport Layer Security (TLS) in 1999
- OpenSSL developed based on public spec
- Mozilla uses NSS (Network Security Services) originally developed by Netscape

Protocol

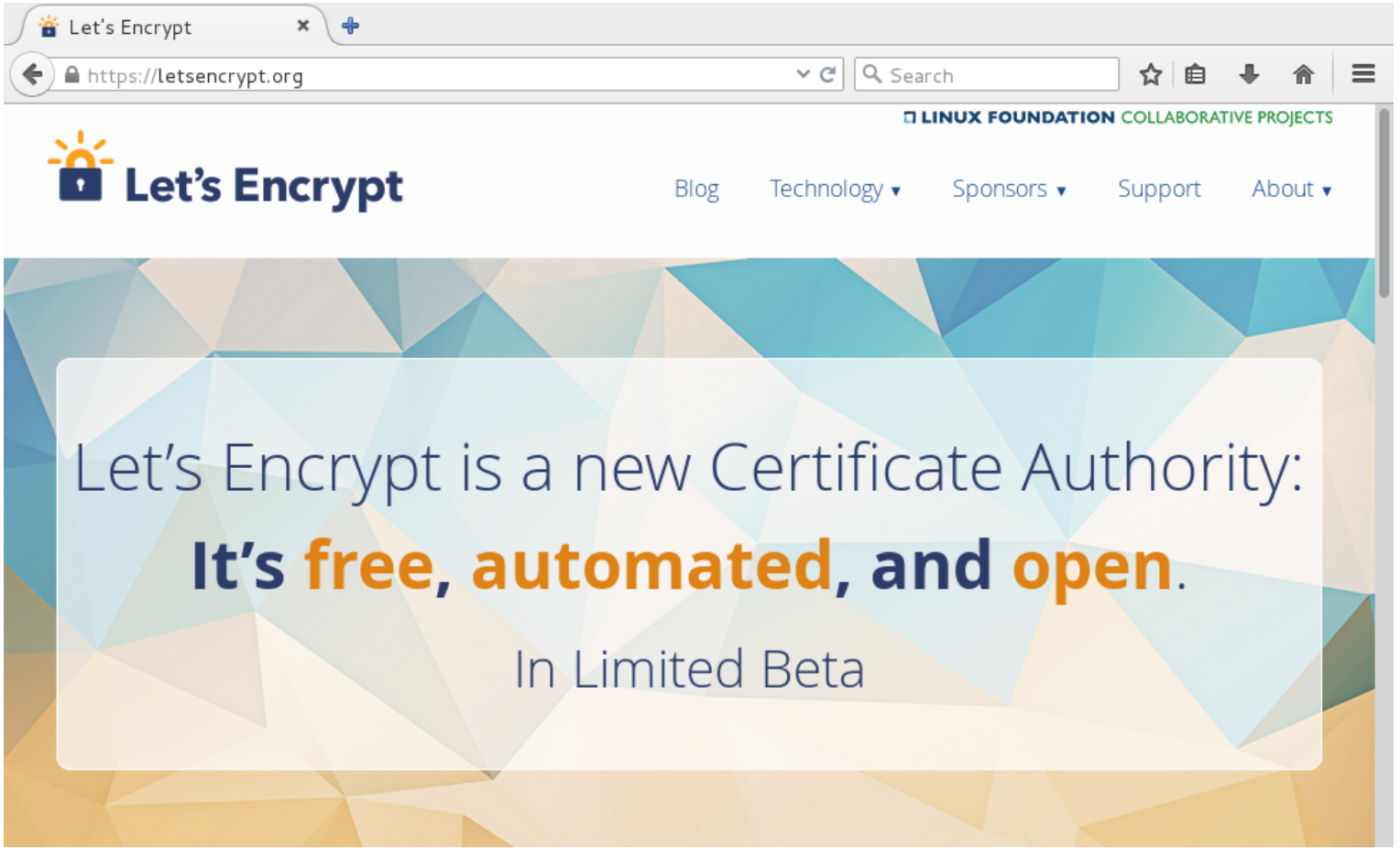

Client
Server
Connect to port 443, send hello & version supported
Sends public key, random number, certificate
Complete DH exchange or encrypt random number

Certificate Authorities

Domain validation

Extended validation

Cost & complexity of certificates



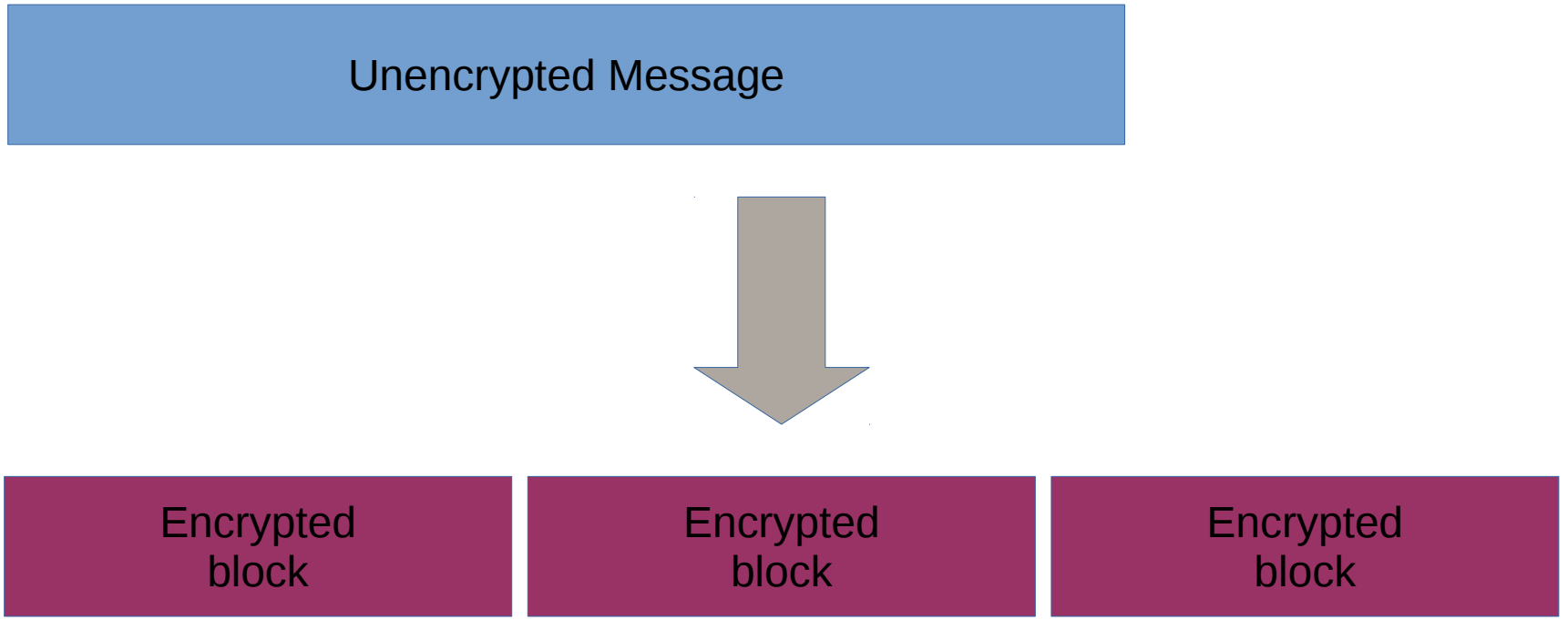
What if the certificate authority is not trustworthy?

- Certificate authorities have issued fraudulent certificates
- Superfish spyware installed its own CA key
- PrivDog allowed any certificate without any validation

Downgrade Attacks

- Old versions of SSL had security weaknesses
- Initial negotiation of supported versions is not encrypted or authenticated
- Active attacker can block initial negotiation and cause fallback to older, insecure version of SSL

Padding Oracle On Downgraded Legacy Encryption (POODLE)



The square-and-multiply algorithm

Calculate: n^2 … n^4 … n^8 … n^{16} … n^{32} …

Example: $n^{42} = n^2 \times n^8 \times n^{32}$

Problem: The square and multiply operations differ in CPU time, power usage, and memory access pattern

- Timing attacks
- Radio frequency emission
- Cache analysis

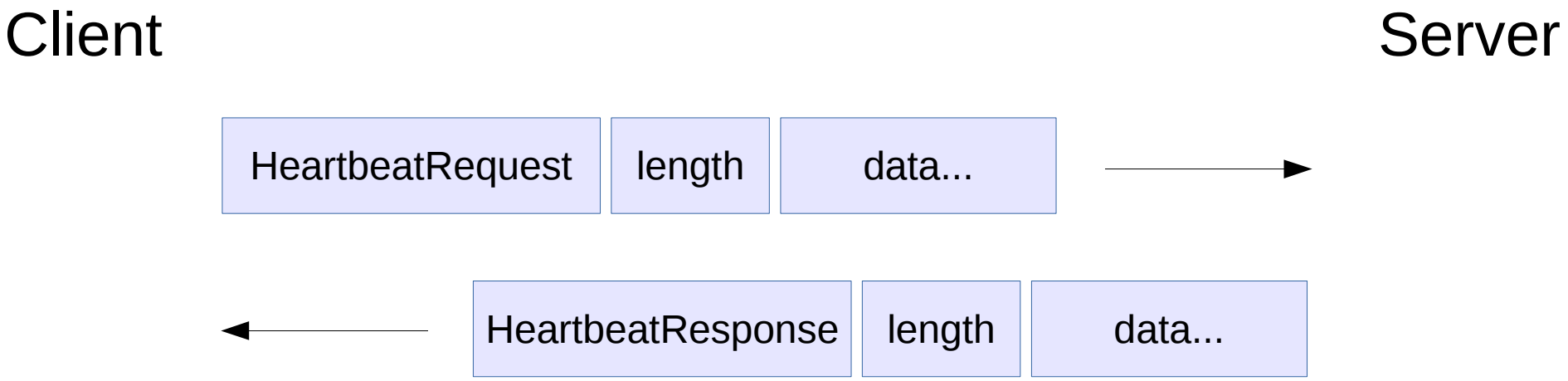
Data Compression Exploits

Even if the encryption is secure, the size of messages may leak information

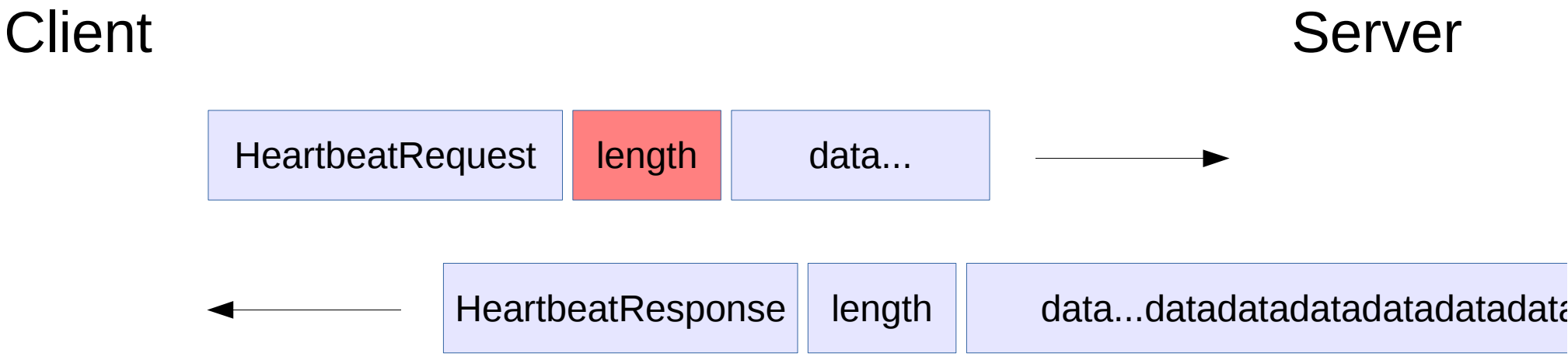
Compression Ratio Info-leak Made Easy (“CRIME")

- Trick client into accessing various URLs (eg embedded images, iframes, javascript)
- Make guesses at session cookie in URL
- If a substring matches, data compression will shorten the message!
- Mitigate by not compressing headers

TLS Heartbeat Extension (RFC 6520)



OpenSSL had a bug...



Heartbleed

- OpenSSL used the requested length rather than the actual length in its response
- Returned whatever happened to be in that area of memory
- Allowed reading up to ~64KB per request
- Potentially leaked private keys

Impact of exposed keys

- Can spoof the server
- If Diffie-Hellman was not used, exposes all prior traffic
- Need to get new certificate for new key
- Need to revoke old certificate

Certificate Revocation

- Even if the certificate is free, revocation often is not (Startcom wanted $25)
- How to communicate revoked status to browsers?

Online Certificate Status Protocol (OCSP)

- Browser checks status of certificate with CA
- Typically unencrypted http
- Generally ignored on timeout
- Responses often lack nonce or expiration

Attacking Diffie-Hellman

The security of Diffie-Hellman key exchange is based on the difficulty of finding discrete logarithms in a finite field.

Logarithms have the property that the sum of logarithms is equal to the logarithm of the product, eg:

$$\log(2) + \log(3) = \log(6)$$

To break Diffie-Hellman, find logarithms of many small primes, then combine them to find the logs of larger numbers.

Avoiding the precomputation attack

- Use a modulus large enough to make this precomputation infeasible (2048+ bits)

  ...but encryption is $O(n^3)$ for n-bit modulus

- Don't use the same modulus as everyone else

  ...but generating strong primes takes time

- Use elliptic curve Diffie-Hellman

Conclusion

- Disable legacy weak crypto
- Don't allow downgrade to SSL 3
- Don't use RC4, DES or other weak algorithms
- Use ephemeral DH with ≥2048-bit modulus, or elliptic curve
- If using finite-field DH, generate a custom strong prime (see weakdh.org)
- Don't mix secure and insecure content